



# Sustainability considerations in container pre-marshalling problem

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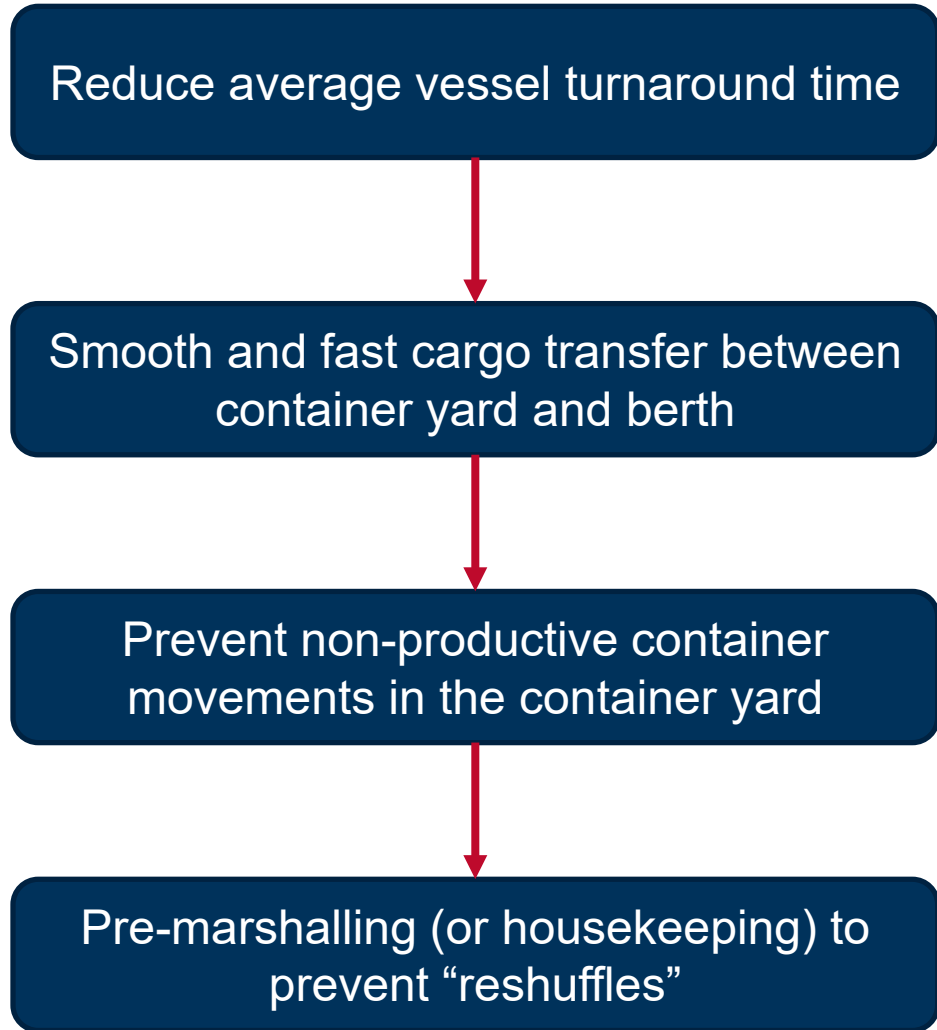
- Research motivations
- Problem definition & literature review
- Model formulation
- Solution heuristic development
- Computational experiments
- Conclusions



(Source: <https://www.termavi.com/en/multimedia-2/>)

A major source of inefficiency  
(Choe et al., 2011)

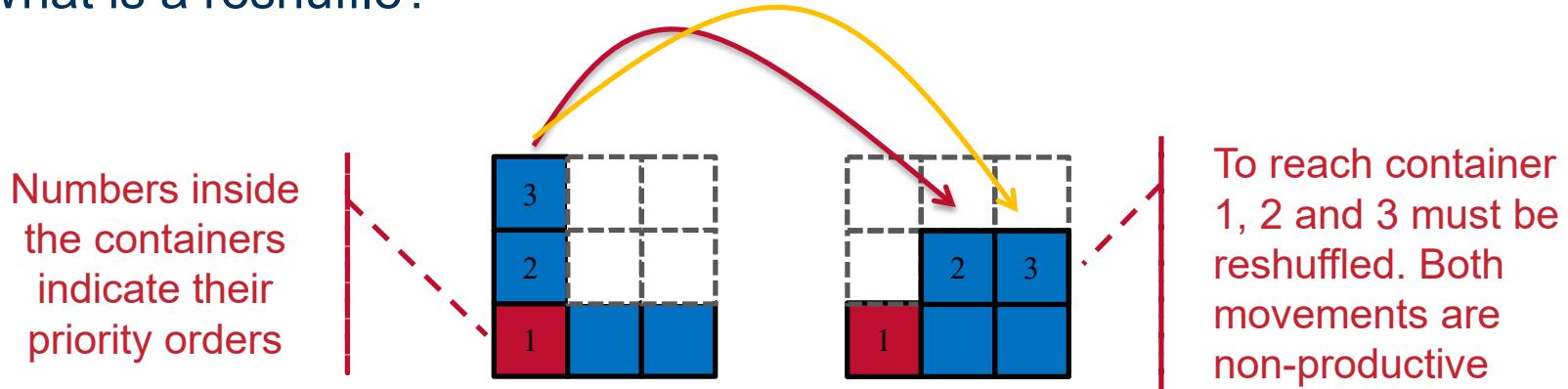
Also called “rehandle”,  
“restow”, or “repositioning”



# Research motivations

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## What is a reshuffle?



## How does pre-marshalling work?

Reorder containers during off-peak hours to eliminate reshuffles during loading/retrieval in a yard bay

## Why is pre-marshalling needed?

Scarcity of empty yard slots  
Random container arrivals  
Imprecise container information

*(Kang et al., 2006; Lee and Chao, 2009; Choe et al., 2011).*

## Environmental sustainability in maritime shipping and port operations

### Common practices

Speed reduction  
Re-routing  
Berth scheduling

*(Mansouri et al., 2015)*

### Main research areas

Emission reduction  
Eco-efficient operations  
Climate change and regulation  
Ship design and mobility  
Carbon footprint case studies

*(Davarzani et al., 2016)*

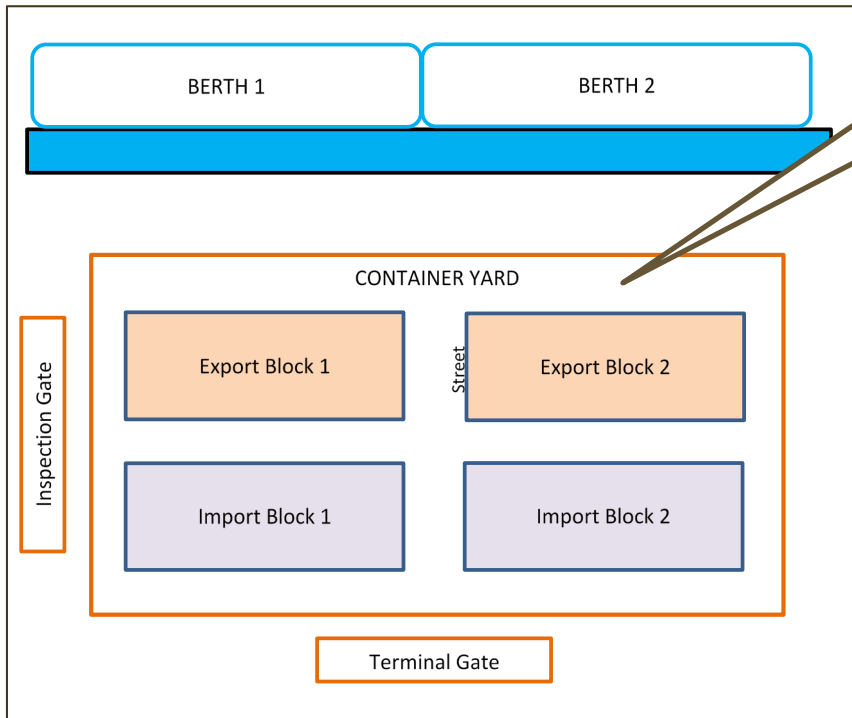
### Regulatory developments

MARPOL Convention (IMO, 1973-1997)  
ISO 14001 (1996, 2004, 2015)  
Ecoports (ESPO, 1997, 2011)

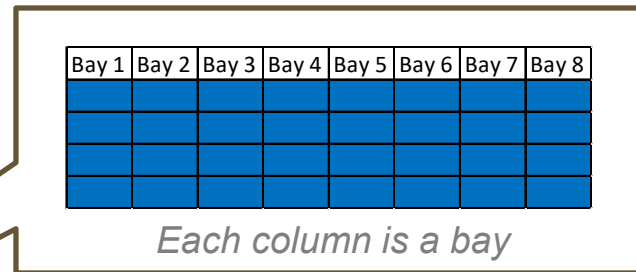
# Problem definition

## Common terminology

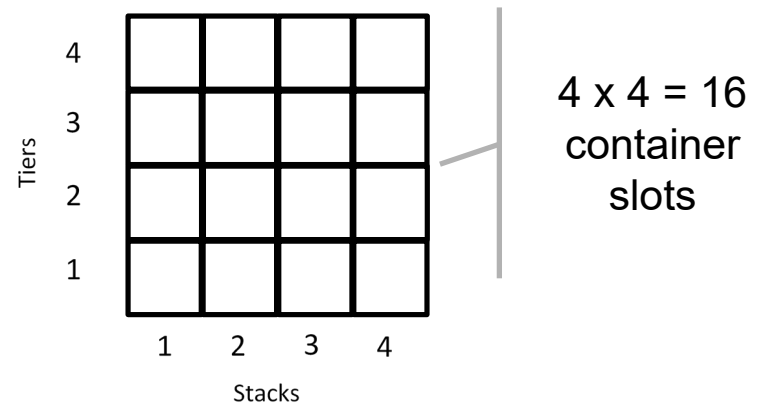
**Blocks (or streets)** are collections of container bays



Container terminal layout



**Bay:** A rectangular configuration of container slots divided column-wise into stacks and row-wise into tiers.

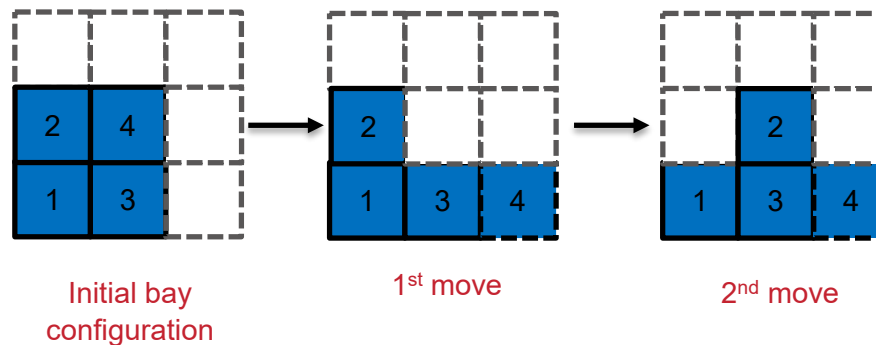


# Problem definition

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A pre-marshalling example

An “unordered” bay configuration: 2 blocks 1, and 4 blocks 3



Move  
container 4  
from stack 2  
to stack 3

An “ordered” bay configuration: None of the containers are blocked by others.

Move  
container 2  
from stack 4  
to stack 1

Container pre-marshalling problem: Minimize **the number of container movements** required to obtain an ordered bay

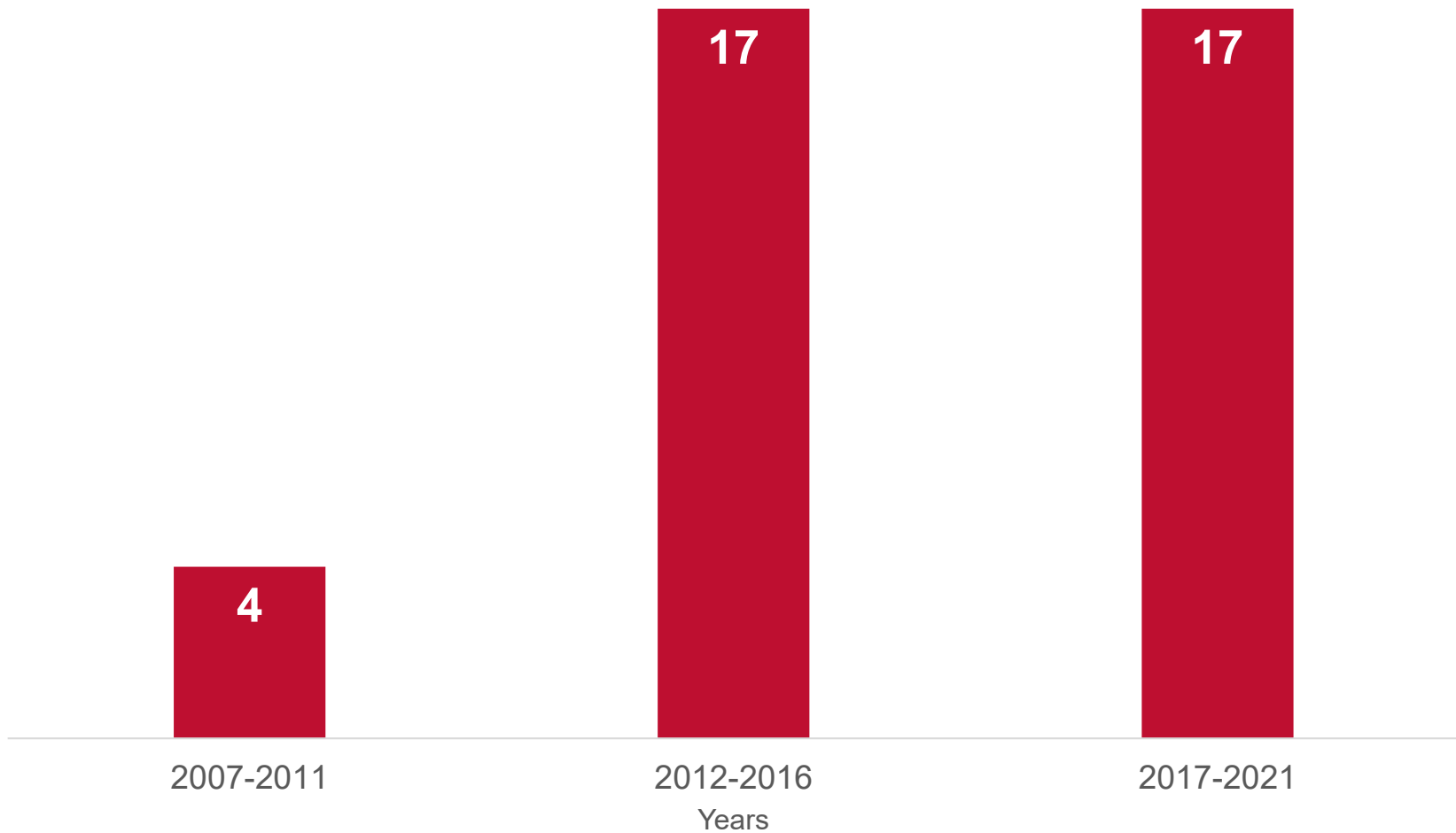
- Lee and Hsu, 2007: Multi-commodity flow formulation
- Lee and Chao, 2009: Neighbourhood search + integer programming
- Bortfeldt and Forster, 2012: Tree search heuristic
- Expósito-Izquierdo et al., 2012: A\* heuristic, instance generator
- Rendl and Prandstetter, 2013: Constraint programming
- Jovanovic et al., 2017: Greedy heuristic
- Tierney et al., 2017: A\* & IDA algorithms
- Tanaka and Tierney, 2018: Iterative deepening branch and bound
- Tierney et al., 2019: Branch and bound algorithm
- Parreño-Torres et al., 2019: New integer programming formulations
- Hottung et al., 2020: Deep learning heuristic tree search



Non-conventional problem definitions appear recently:

- Parreño-Torres et al., 2020: Minimize crane times in pre-marshalling problems
- Boge et al., 2020: Robust optimization for premarshalling with uncertain priority classes
- Zweers et al., 2020: Stochastic container relocation problem with constrained pre-processing

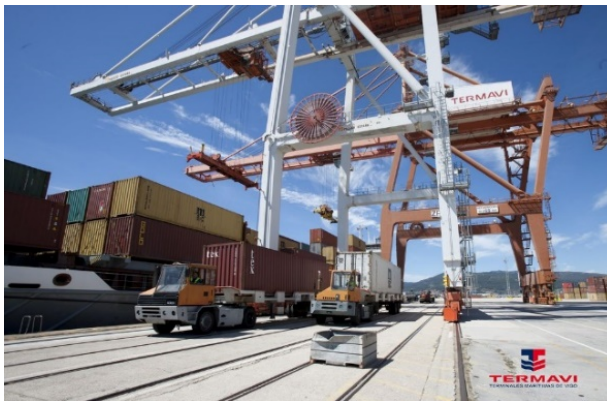
Number of publications





## Container pre-marshalling problem with sustainability considerations

For a given initial bay configuration and the (retrieval) priority order of the containers, determine the sequence of the container movements so as to minimize **the total energy consumption of the yard cranes** and achieve an **ordered bay**.



(Source: <https://www.termavi.com/en/multimedia-2/>)

## Assumptions:

- All containers are assumed to be of the same size and dimensions (1 TEU).
- The priority order of the containers for retrieval is known in advance.
- Only one crane operates on the bay, and only a single container can be moved at a time.
- The number of containers in the bay does not change during the pre-marshalling operation.
- All stacks and tiers in the bay are identical.
- Each slot can contain at most one container.
- The unit idle and operational energy consumption of the cranes are known and constant for each container.

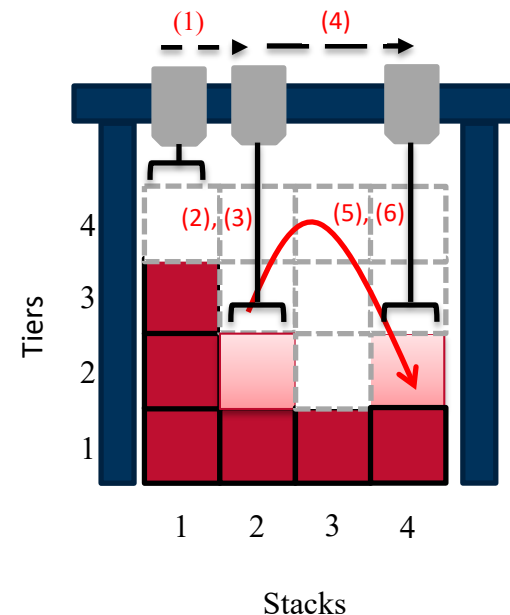
*(Lee and Hsu, 2007; Lee and Chao, 2009; Wang et al., 2015; 2017; Tierney et al., 2017; Tanaka and Tierney, 2018)*

# Model formulation

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- The proposed mixed integer programming model is formulated by partially adopting the classical pre-marshalling problem formulation proposed by Parreño-Torres et al. (2019).
- The crane energy consumption of a single container movement is calculated by considering six crane manoeuvres:

1. Movement of the crane's trolley from its previous position to the source stack  $s$
2. Lowering of the spreader to the topmost container of stack  $s$
3. Hoisting of the topmost container with the spreader
4. Movement of the crane's trolley from  $s$  to the target stack  $k$
5. Lowering of the spreader to place the container on top of stack  $k$
6. Hoisting of the spreader back to the trolley



<b>Set</b>	<b>Explanation</b>
$\mathcal{S}$	Set of stacks, $s, k, l = \{1, \dots, S\}$
$\mathcal{H}$	Set of tiers, $h, e = \{1, \dots, H\}$
$\mathcal{T}$	Set of time segments, $t = \{1, \dots, T\}$
$\mathcal{P}$	Set of container retrieval priorities, $p, r = \{1, \dots, P\}$

<b>Parameter</b>	<b>Explanation</b>
$N_p$	The number of containers in priority group $p$ ,
$tfl_{sk}$	The energy consumed for moving the trolley loaded with a container from stack $s$ to $k$
$tfu_{ls}$	The energy consumed for moving the trolley unloaded from stack $l$ to $s$
$tfi_h$	The energy consumed by keeping the trolley idle during the spreader's movement between trolley level and tier $h$
$sfl_h$	The energy consumed for lowering the spreader empty to tier $h$ and hoisting a container to the trolley level
$sfu_h$	The energy consumed for lowering the spreader loaded with a container to tier $h$ and hoisting it empty to the trolley level
$sfi_{sk}$	The energy consumed by keeping the spreader idle during the trolley's movement from stack $s$ to $k$

## Decision variables

- $Y_{shkept}$  1 if a container of priority  $p$  is moved from stack  $s$  tier  $h$  to stack  $k$  tier  $e$  between the time segments  $t$  and  $t + 1$ , 0 otherwise  
 $\forall s \in \mathcal{S}, h \in \mathcal{H}, k \in \mathcal{S} \setminus \{s\}, e \in \mathcal{H}, p \in \mathcal{P}, t \in \mathcal{T} \setminus \{T\}$
- $X_{shpt}$  1 if a container with priority  $p$  occupies stack  $s$  tier  $h$  at the beginning of time segment  $t$ , 0 otherwise  $\forall s \in \mathcal{S}, h \in \mathcal{H}, p \in \mathcal{P}, t \in \mathcal{T}$
- $C_{lst}$  1 if the trolley is moved from stack  $l$  to  $s$  for moving a container from stack  $s$  between the time segments  $t$  and  $t + 1$ , 0 otherwise  $\forall l \in \mathcal{S}, s \in \mathcal{S} \setminus \{l\}, t \in \mathcal{T} \setminus \{1, T\}$

# Model formulation

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Minimise

$$\sum_{l=1}^S \sum_{s=1}^S \sum_{t=2}^{T-1} (tfu_{ls} * C_{lst}) +$$

(1) Minimize the total energy consumption of the crane to perform all movements

(1)

$$\sum_{s=1}^S \sum_{h=1}^H \sum_{\substack{k=1 \\ k \neq s}}^S \sum_{e=1}^H \sum_{p=1}^P \sum_{t=1}^{T-1} \left( (tfl_{sk} + tfi_h + tfi_e + sfi_{sk} + sfl_h + sfu_e) * Y_{shkept} \right)$$

such that

$$\sum_{s=1}^S \sum_{h=1}^H X_{shpt} = N_p \quad \forall p \in \mathcal{P}, t \in \mathcal{T}$$

(2) Assign a slot position to all containers at all time segments

(2)

$$\sum_{s=1}^S \sum_{h=1}^H \sum_{\substack{k=1 \\ k \neq s}}^S \sum_{e=1}^H \sum_{p=1}^P Y_{shkept} \leq 1 \quad \forall t \in \mathcal{T} \setminus \{T\}$$

(3) 1 movement per time segment

(3)

$$\sum_{s=1}^S \sum_{h=1}^H \sum_{\substack{k=1 \\ k \neq s}}^S \sum_{e=1}^H \sum_{p=1}^P Y_{shkept+1} \leq \sum_{s=1}^S \sum_{h=1}^H \sum_{\substack{k=1 \\ k \neq s}}^S \sum_{e=1}^H \sum_{p=1}^P Y_{shkept} \quad \forall t \in \mathcal{T} \setminus \{T-1, T\}$$

(4) Perform the movements at earliest time segments

(4)



# Model formulation

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$$\sum_{p=1}^P X_{s1pt} \leq 1 \quad \forall s \in \mathcal{S}, t \in \mathcal{T} \setminus \{1\} \quad (5)$$

(5)-(6) Each tier at most 1 container

$$\sum_{p=1}^P X_{sh+1pt} \leq \sum_{p=1}^P X_{shpt} \quad \forall s \in \mathcal{S}, h \in \mathcal{H} \setminus \{H\}, t \in \mathcal{T} \setminus \{1, T\} \quad (6)$$

(7) Only move the topmost container at each stack

$$X_{shpt} + \sum_{r=1}^P X_{sh+1rt} \leq 1 + X_{shpt+1} \quad \forall s \in \mathcal{S}, h \in \mathcal{H} \setminus \{H\}, p \in \mathcal{P}, t \in \mathcal{T} \setminus \{T\} \quad (7)$$

$$\sum_{\substack{s=1 \\ s \neq k}}^S \sum_{h=1}^H \sum_{e=1}^H Y_{shkept} + \sum_{\substack{s=1 \\ s \neq k}}^S \sum_{h=1}^H \sum_{e=1}^H Y_{keshpt+1} \leq 1 \quad \forall k \in \mathcal{S}, p \in \mathcal{P}, t \in \mathcal{T} \quad (8)$$

(8) Prevent transitive moves (Tierney et al., 2017)

$$\sum_{\substack{k=1 \\ k \neq s}}^S \sum_{e=1}^H Y_{shkept} + \sum_{\substack{k=1 \\ k \neq s}}^S \sum_{e=1}^H Y_{keshpt+1} \leq 1 \quad \forall s \in \mathcal{S}, h \in \mathcal{H}, p \in \mathcal{P}, t \in \mathcal{T} \setminus \{T-1, T\} \quad (9)$$

(9) Prevent symmetrical moves (Tanaka and Tierney, 2018)

$$X_{shpt} + \sum_{\substack{k=1 \\ k \neq s}}^S \sum_{e=1}^H Y_{keshpt} = X_{shpt+1} + \sum_{\substack{k=1 \\ k \neq s}}^S \sum_{e=1}^H Y_{shkept} \quad \forall s \in \mathcal{S}, h \in \mathcal{H}, p \in \mathcal{P}, t \in \mathcal{T} \setminus \{T\} \quad (10)$$

(10) Slot occupation constraint

# Model formulation

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$$\sum_{r=p}^P X_{sh+1rT} \leq \sum_{r=p}^P X_{shrT} \quad \forall s \in \mathcal{S}, h \in \mathcal{H} \setminus \{H\}, p \in \mathcal{P}$$

(11) The bay is ordered at T

$$C_{lst} \leq \sum_{h=1}^H \sum_{\substack{k=1 \\ k \neq s}}^S \sum_{e=1}^H \sum_{p=1}^P Y_{shkept} \quad \forall l \in \mathcal{S}, s \in \mathcal{S}, t \in \mathcal{T} \setminus \{1, T\}$$

(12)-(14)  $C_{lst}$  is 1 if and only if the target stack of the movement during  $t-1$  was  $l$  and the source stack of the movement during  $t$  is  $s$

$$C_{lst} \leq \sum_{\substack{k=1 \\ k \neq l}}^S \sum_{e=1}^H \sum_{h=1}^H \sum_{p=1}^P Y_{kelhpt-1} \quad \forall l \in \mathcal{S}, s \in \mathcal{S}, t \in \mathcal{T} \setminus \{1, T\}$$

(13)

$$\sum_{\substack{k=1 \\ k \neq l}}^S \sum_{e=1}^H \sum_{h=1}^H \sum_{p=1}^P Y_{kelhpt-1} + \sum_{h=1}^H \sum_{\substack{k=1 \\ k \neq s}}^S \sum_{e=1}^H \sum_{p=1}^P Y_{shkept} \leq C_{lst} + 1 \quad \forall l \in \mathcal{S}, s \in \mathcal{S}, t \in \mathcal{T} \setminus \{1, T\}$$

(14)

$$Y_{shkept}, C_{lst} \in \mathbb{R}^+, X_{shpt} \in \{0,1\} \quad \forall s, k, l, e, h, p, t$$

(15) Integrality and nonnegativity constraints

# Solution heuristic development

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- An A\* algorithm-based solution heuristic is being developed.
- A\* algorithm is a well-known pathfinding algorithm and was proposed to solve the conventional pre-marshalling problem (Exposito-Izquierdo et al., 2012; Tierney et al., 2017)
- In a graph of **nodes**, iteratively construct the shortest path between a **starting location** and a **destination** using a function on the sum of **exact and estimated costs** (distances):

$$f(n) = g(n) + h(n)$$

The exact cost  
from start to  
node  $n$

The estimated  
cost from node  
 $n$  to destination

# Solution heuristic development

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- In the proposed solution method, the energy consumption function is incorporated into the A\* algorithm, and path selection is made based on the realized and estimated energy consumptions.
- Assume initial bay configuration as the starting location and any ordered bay configuration as the destination.
- Each intermediate bay configuration refers to the nodes visited along the path between the starting location and destination.
- The algorithm favours the paths which have the lowest total energy consumption.
- (Under development): To intensify the search in promising paths, hybridize the A\* algorithm with an evolutionary heuristic.
  - Select “ $m$ ” elite paths with minimum total energy consumption
  - Generate random move sequences from each elite path.
  - Apply evolutionary mechanisms to optimize the move sequences
  - Stop if a sequence yields an ordered bay.

# Computational experiments

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Experiment instances:

	Stacks	Tiers	Containers
Lee and Hsu (2007)	6	4	14
Expósito-Izquierdo et al. (2012)	4	4	8 / 12 / 16
	7	4	7 / 14 / 28
	10	4	10 / 20 / 40
This paper	6	4	10
Use case Spanish terminal	6	4	Various

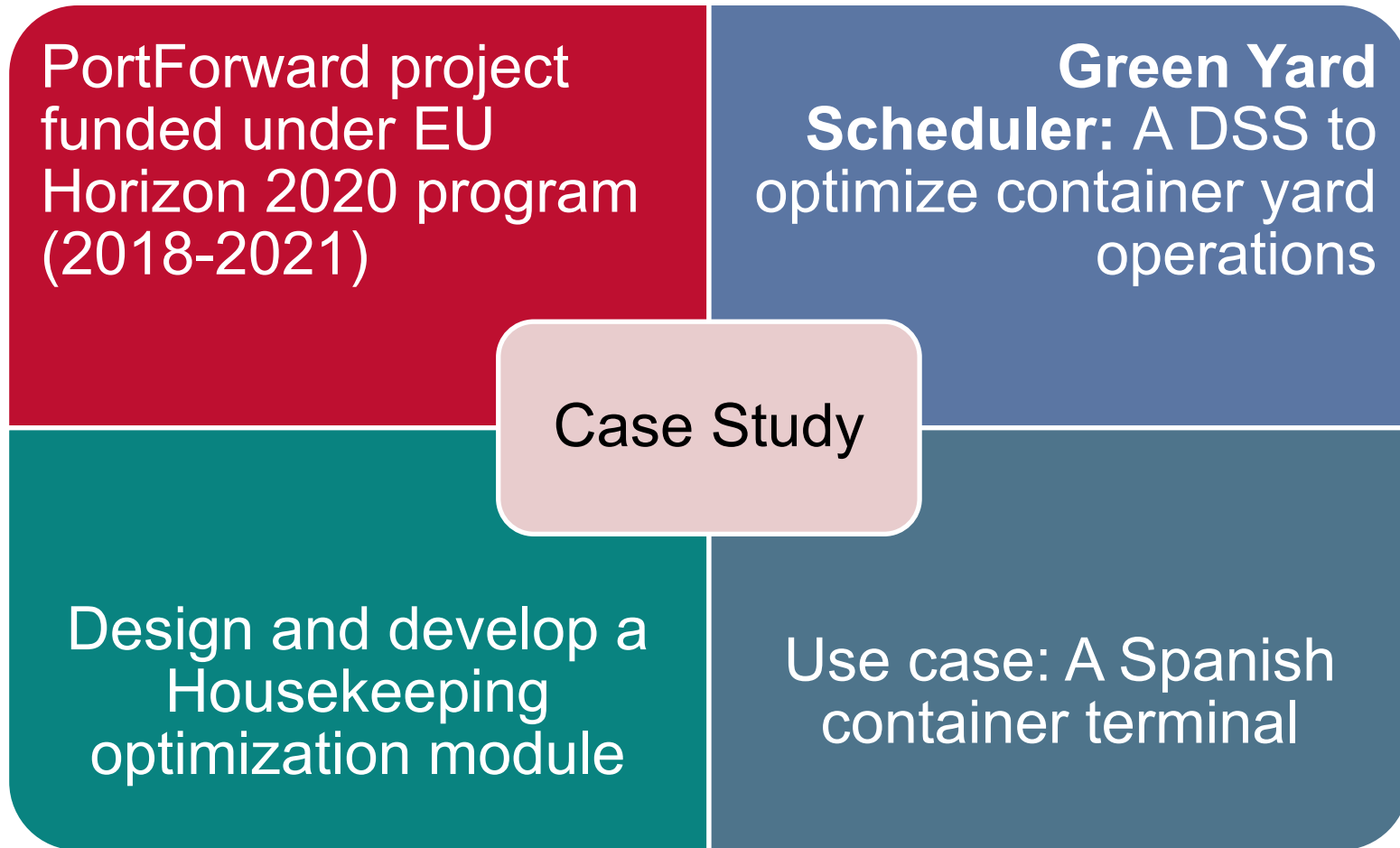
## Preliminary observations

Container movements and energy consumption do not conflict with each other according to the observations

**4-6%** energy reduction with the proposed objective function

The solution heuristic can find good solutions for instances with low bay occupation ratio (=containers/total bay capacity) in a reasonable computation time.

The heuristic slows down with difficult instances (i.e. high bay occupation ratio and/or a bay with more stacks/tiers) due to the memory issues.





# Conclusions

- Container pre-marshalling problem with sustainability considerations is introduced.
- An integer programming model formulation is presented and a solution heuristic is discussed.
- Preliminary experiment results are discussed.

## Future research:

- Complete the hybridization of A\* algorithm with evolutionary heuristic to find good solution for difficult pre-marshalling problem instances.
- Expand the research to intra-bay pre-marshalling movements.

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# Thank you for your attention!

## Questions, comments?

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